

## **Compiler Construction WS11/12**

# **Exercise Sheet 5**

Please hand in the solutions to the theoretical exercises until the beginning of the lecture next Friday 2011-11-25, 12:00. Please write the number of your tutorial group or the name of your tutor on the first sheet of your solution.

### Exercise 5.1 LR(0) (Points: 6+2+2+2+4)

Let the grammar  $G = (\{S', S, A, B, C\}, \{a, b, c, d\}, P, S')$  with productions P:

$$\begin{array}{rcl} S' & \rightarrow & S \\ S & \rightarrow & A \, B \mid A \\ A & \rightarrow & a \, C \, c \\ C & \rightarrow & b \, b \, C \mid b \\ B & \rightarrow & c \, d \end{array}$$

- 1. Construct the  $LR_0(G)$  automaton with the direct construction algorithm from the lecture.
- 2. Mark all inadequate states in the  $LR_0(G)$  automaton. For each inadequate state you have to enumerate all the conflicts (each conflict is a pair of items) and classify them.
- 3. Is G an SLR(1) grammar? Justify your answer.
- 4. Construct  $LR_1(G)$  by adding lookahead sets. To keep your write-up short, only construct the LR(1)-items for the conflicting items in the LR(0)-inadequate states.
- 5. Give a successful run of the PDA  $P_1(G)$  controlled by  $LR_1(G)$  on the input word w = a b b b b b c c d. You can do this by creating a table containing columns for the current stack content, the remaining input and the next action. You do not need to formally specify  $P_1(G)$ . At which points of the run would there be conflicts if it was not for the lookahead sets added and why does your selection of the lookahead sets prevent these situations?

### Exercise 5.2 LL(0) and LR(0) (Points: 2+2+2)

Prove or disprove the following claims:

- 1. All LL(0) languages are also LR(0) languages.
- 2. All regular languages are LR(0).
- 3. Not all LR(0) languages are regular.

#### **Exercise 5.3 Project**

Next project assignment is available online.